

Today

Multiplicative Weight
Update Method
(part V)

- Hardcore set lemma
- XOR Lemma
- Wrapup

CSS.205.1

Toolkit in TCS

- Lecture # 19

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Application: Impagliazzo's Hardcore Set
Lemma

Previously, proved using von-Neumann's
Minimax Theorem

Today, "Constructive" proof using MWUM.

(Hardcore Set Lemma \Rightarrow XOR Lemma)

Hard-core Sets:

Boolean functions: $f: \{0,1\}^n \rightarrow \{0,1\}$
Boolean hypercube Boolean
(replaced by any set X)

Model of computation: Circuits.
(DAG of internal nodes)

corresponding to binary 1,
binary 0 or unary NOT gates
& leaves - input variables)

Size of such circuit = #gates of the circuit.

How well does a ckt of size at most S compute $f: \{0,1\}^n \rightarrow \{0,1\}$?

$$\delta(C, f) \triangleq \Pr_{x \in \{0,1\}^n} [C(x) = f(x)].$$

① Worst case Hardness:

$$\delta(C, f) < 1 \text{ for all ckt's of size } S.$$

② Average-case Hardness

(a) Mildly average-case hard.

ϵ -weakly hard ($\epsilon \in (0, 1/2)$)
against ckt's of size S .

if \forall ckt's C of size (at most) S

$$\delta(C, f) \leq 1 - \epsilon.$$

(b) Strongly average-case hard

γ -strongly hard against ckt of size S

($\gamma \in (0, 1/2)$)

if \forall ckt C of size (at most) S

$$\delta(C, f) \leq \frac{1}{2} + \gamma$$

→ Yao's XOR Lemma: Mildly average case hard f

⇔ Strongly average-case hard f'

$f: \{0, 1\}^n \rightarrow \{0, 1\}$

$f' = f^{\oplus k}: \{0, 1\}^{nk} \rightarrow \{0, 1\}$

$(x_1 \dots x_k) \mapsto f(x_1) \oplus f(x_2) \oplus \dots \oplus f(x_k)$

Yao's XOR Lemma:

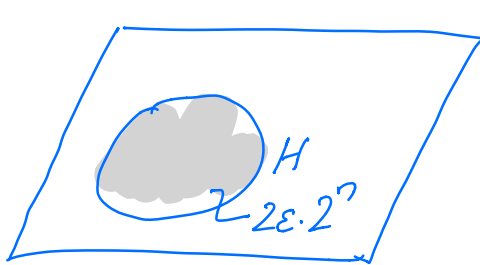
f is ϵ -weakly hard against ckt of size S

⇔

$f^{\oplus k}$ is $\gamma + (1-\epsilon)^k$ -strongly hard against ckt of size $S' = O(\epsilon^2 \gamma^2 S)$

Proof (due to Impagliazzo) via Hardcore Sets

Hardcore Sets



$\{0,1\}^n$

$$f: \{0,1\}^n \rightarrow \{0,1\}$$

f is ϵ -weakly hard.
against ckt's of
size S

\forall ckt's C of size S

$$\Pr_{x \leftarrow \{0,1\}^n} [f(x) = C(x)] \leq 1 - \epsilon$$

$H \subseteq \{0,1\}^n$ is r -hardcore set for f
against ckt's of size S .

if \forall ckt's C of size S $\Pr_{x \leftarrow H} [C(x) = f(x)] \leq \frac{1}{2} + \delta$

Obs: $H \subseteq \{0,1\}^n$. $|H| = 2\epsilon \cdot 2^n$ is a
 r -hardcore set for ckt's of size S

\Downarrow
 f is $(\epsilon-r)$ -weakly hard against ckt's of
size S .

Hardcore set Lemma: Converse to
the observation.

H - set of size $\Theta(\epsilon)$.

H - hardcore distribution
 ϵ -smooth distribution.

For any $\epsilon \in (0, 1)$, $\mathcal{D} \sim \{0, 1\}^n$ is said to be ϵ -smooth

$$\text{if } \forall x \in \{0, 1\}^n, \Pr_{X \sim \mathcal{D}} [X=x] \leq \frac{1}{\epsilon \cdot 2^n}$$

eg: ϵ -smooth distribution.

(1) uniform dist on $\{0, 1\}^n$

(2) $H \subseteq \{0, 1\}^n$. $|H| = \epsilon \cdot 2^n$ } ϵ -smooth
 \mathcal{D} - unif dist on H } ϵ -flat distrib.

(3) ϵ -smooth dist is a convex combination of ϵ -flat dist
(we won't use this)

(4) $\mathcal{P} = \{\mathcal{D} \mid \mathcal{D} \text{ is } \epsilon\text{-smooth}\}$

$$\mathcal{D}: \{0, 1\}^n \rightarrow [0, 1]$$

\mathcal{P} - convex set.

Impagliazzo's Hardcore Set Lemma:

$f: \{0, 1\}^n \rightarrow \{0, 1\}$ is ϵ -weakly hard against
cts of size S

$\forall r \in (0, 1/2)$
 then there is a ϵ -smooth dist H
 s.t. f is r -strongly hard against ckt's of
 size $S' = O\left(\frac{r^2 S}{\log(\frac{1}{\epsilon})}\right)$ on H

$\forall (r, \forall$ ckt's C of size S'
 $\Pr_{x \sim H} [C(x) = f(x)] \leq \frac{1}{2} + r$

Pf: By contradiction

Suppose for every ϵ -smooth distribution H
 there is a ckt C of size S'

s.t. $\Pr_{x \sim H} [f(x) = C(x)] \geq \frac{1}{2} + r.$

(Last time: weak learner against every dist

\Downarrow
 strong learner)

Now: weak learner against ϵ -smooth dist

\Downarrow
 strong learner.

Catch: $P^{(f)}$ - MWUM outputs must
 be an ϵ -smooth dist.)

Boosting:

1. Initialize $P^{(1)} \leftarrow$ Unit dist on $\{x_i\}^T$
2. For $t \leftarrow 1$ to T .

(a). Construct the ckt C of size $\leq S'$ that

$$\Pr_{x \leftarrow P^{(t)}} [C(x) = f(x)] \geq \frac{1}{2} + \gamma$$

(b) Update

$$\tilde{P}^{(t+1)}(x) \leftarrow P^{(t)}(x) \frac{1 - \eta m_t^{(t)}(x)}{\int P^{(t)}(x) (1 - \eta m_t^{(t)}(x)) dx}$$

$\underbrace{\tilde{P}^{(t+1)}}_{\epsilon\text{-smooth}} \quad m_t^{(t)}(x) = 1 - |C(x) - f(x)|$

Project $\tilde{P}^{(t+1)}$ to \mathcal{P} to obtain an ϵ -smooth

$$P^{(t+1)} = \min_{P \in \mathcal{P}} RE(\tilde{P}^{(t+1)} \| P) \text{ distribution}$$

MWUM (rel entropy).

$$\sum_{t=1}^T \eta m_t^{(t)} \cdot P^{(t)} \leq (1 + \eta) \sum \eta m_t^{(t)} \cdot P + RE(P \| P^{(1)})$$

$\forall P \in \mathcal{P}.$

MWUM: Majority of T ckt's $C_1 \dots C_T$ is an $(1 - \epsilon)$ -approximation to f

$$\text{If } T = \left\lceil \frac{2}{\epsilon^2} \log\left(\frac{1}{\epsilon}\right) \right\rceil$$

$$S' = O\left(\frac{\epsilon^2 S}{\log\left(\frac{1}{\epsilon}\right)}\right) \quad \square$$

Proof of Yao's XOR Lemma.

Suppos

$f: \{0,1\}^n \rightarrow \{0,1\}$ is ϵ -weakly hard against
 ccts of size S .

$$\text{Pr}_{x \leftarrow \{0,1\}^n} [C(x) = f(x)] \leq 1 - \epsilon, \quad \forall C \text{ of size } \leq S$$

\Downarrow Hardcore Set Lemma.

$$\forall r \in (0, 1/2)$$

\exists ϵ -hardcore dist H on $\{0,1\}^n$ s.t

$$\text{Pr}_{x \leftarrow H} [C(x) = f(x)] \leq \frac{1}{2} + r \quad \forall C \text{ of size}$$

$$O\left(\frac{r^2 S}{\log\left(\frac{1}{\epsilon}\right)}\right)$$

\Downarrow Proof of XOR Lemma.

(Needs to be proved).

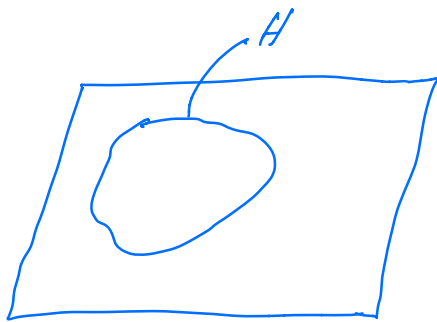
$$\forall r \in (0, 1/2).$$

$$f^{\oplus k}: \{0,1\}^{nk} \rightarrow \{0,1\}$$

$$(x_1, \dots, x_k) \mapsto \bigoplus_{i=1}^k f(x_i)$$

is $r + (1-\epsilon)^k$
 -strongly
 hard against

cuts of size $O\left(\frac{r^5}{\log(\frac{1}{\epsilon})}\right)$



$\{0,1\}^n$

U_n - uniform dist on $\{0,1\}^n$

H - ϵ -smooth distribution

$$\forall x, H(x) \leq \frac{1}{\epsilon \cdot 2^n}$$

Define distribution G on $\{0,1\}^n$

$$G(x) = \frac{\frac{1}{2^n} - \epsilon H(x)}{1 - \epsilon} = \frac{U_n(x) - \epsilon H(x)}{1 - \epsilon}$$

i.e., $U_n(x) = \epsilon \cdot H(x) + (1 - \epsilon) G(x)$.

Qns (1) $G(x) \geq 0 \quad \forall x$?

(equiv to $H(x) \leq \frac{1}{\epsilon \cdot 2^n}$)

(2) $\sum_x G(x) = 1$.

(since $\sum H(x) = 1 = \sum U_n(x) = 1$)

} - distribution

Hence, G is valid distribution on $\{0,1\}^n$

$U_n(x) = \epsilon H(x) + (1 - \epsilon) G(x), \quad \forall x \in \{0,1\}^n$

U - $(\epsilon, 1 - \epsilon)$ convex combination of the 2 dist H & G .

U - can be generated as follows:

- first toss a coin $P_H[\text{heads}] = \epsilon$
 $P_H[\text{tails}] = 1 - \epsilon$
- If coin = heads, then op a sample from H
- otherwise op a sample from G .

How well do cks compare $f(x_1) \oplus f(x_2) \oplus \dots \oplus f(x_k)$?

$k=2$:

$$P_{x_1, x_2 \leftarrow U_n \times U_n} [C(x_1, x_2) = f(x_1) \oplus f(x_2)]$$

For any 2 dist D_1, D_2

$$P_{x_1, x_2 \leftarrow D_1 \times D_2} [C(x_1, x_2) = f(x_1) \oplus f(x_2)] = P_{D_1, D_2}$$

Suppose for contradiction

$$\text{assume } P_{U_1, U_2} \geq \frac{1}{2} + r + (1-\epsilon)^2$$

$$\frac{1}{2} + r + (1-\epsilon)^2 \leq P_{U_1, U_2}$$

$$U_1 = \epsilon H_1 + (1-\epsilon) G_1$$

$$U_2 = \epsilon H_2 + (1-\epsilon) G_2$$

$$\begin{aligned}
 P_{G_1, G_2} &= P_{\varepsilon H_1 + (1-\varepsilon)G_1, \varepsilon H_2 + (1-\varepsilon)G_2} \\
 &= \varepsilon^2 P_{H_1, H_2} + \varepsilon(1-\varepsilon) P_{H_1, G_2} + (1-\varepsilon)\varepsilon P_{G_1, H_2} \\
 &\quad + (1-\varepsilon)^2 P_{G_1, G_2}
 \end{aligned}$$

} U can be simulated using $H_2 \rightarrow G$

$$P_{G_1, G_2} \leq 1$$

$$\frac{1}{2} + \gamma + (1-\varepsilon)^2 \leq \varepsilon^2 P_{H_1, H_2} + \varepsilon(1-\varepsilon) P_{H_1, G_2} + P_{G_1, H_2} \varepsilon(1-\varepsilon) + (1-\varepsilon) P_{G_1, G_2}^2$$

$$\frac{1}{2} + \gamma \leq \varepsilon^2 P_{H_1, H_2} + \varepsilon(1-\varepsilon) P_{H_1, G_2} + \varepsilon(1-\varepsilon) P_{G_1, H_2}$$

At least one P_{H_1, H_2} , P_{H_1, G_2} or $P_{G_1, H_2} \geq \frac{1}{2} + \gamma$

Assume $P_{H_1, G_2} \geq \frac{1}{2} + \gamma$.

$$\frac{1}{2} + \gamma \leq P_{x_1} \left[C(x_1, x_2) = f(x_1) \oplus f(x_2) \right]$$

$x_1 \leftarrow H$
 $x_2 \leftarrow G$

$\exists x_2 \sim G$ st above is true.

Fix that x_2 .

$$\frac{1}{2} + \epsilon \leq P_{x \leftarrow H} [C(x_1, x_2) \oplus f(x_2) = f(x)]$$

That means the ckt $C(x_1, x_2) \oplus f(x_2)$
 compute f correctly on H w/p $\frac{1}{2} + \epsilon$
 (but this contradicts Hardcore set Lemma).

Over assumption that $C(x_1, x_2)$
 computes $f(x_1) \oplus f(x_2)$ w/p $\geq \frac{1}{2} + \epsilon + (1-\epsilon)^t$
 is false

