

Lecture 8 (Guest Lecturer: Yuval Filmus)

Lovász & Functions & its Applications (Notes by Prahladh)

Pronunciation Guide

Hungarian:	S,	SZ,	ZS
Polish	S	SZ	
	Z	Z	

Traffic Light Puzzle:

1 traffic light - controlled by n 3-way switches.
Light changes if all switches changes.

Given a graph $G = (V, E)$

Goal. Find an upper bd on independent set of G .

Theorem (Lovász)

M -symmetric $V \times V$ matrix st

$M(x, y) = 1$ whenever (x, y) is not an edge

$$\Downarrow \\ \alpha(G) \leq \lambda_{\max}(M)$$

Proof: F -independent set, $\mathbb{1}_F$ -char vector

$$\mathbb{1}_F^T M \mathbb{1}_F = \sum_{x, y \in F} M(x, y) = |F|^2$$

On the other hand

$$\mathbb{1}_F^T M \mathbb{1}_F \leq \lambda_{\max} \mathbb{1}_F^T \mathbb{1}_F = \lambda_{\max} |F|$$

Hence, $|F| \leq \lambda_{\max}$



Denote by $\theta(G)$ the best bound that the lemma can get

$$\theta(G) = \min \lambda_{\max}(M) \quad \begin{cases} M(x,y) = 1 & (x,y) \notin E \\ M \leq \lambda_{\max} \end{cases}$$

Erős-Ko-Rado Theorem:

Suppose $k \leq n/2$, if $\mathcal{F} \subseteq \binom{[n]}{k}$ is intersecting then (1) $|\mathcal{F}| \leq \binom{n-1}{k-1}$

(2) $k < n/2$ & $|\mathcal{F}| = \binom{n-1}{k-1}$, then $\mathcal{F}$ is a star.

Proof: Construct $G = (V, E)$ (using Hoffman bd.)

$V = \binom{[n]}{k}$, $A, B \in V$, $A \sim B$ if $A \cap B = \emptyset$
Kneser graph. $(K_n(n, k))$

(Using M adjacency matrix $K_n(n, k)$ yields (1))

Lovász Theorem (Strengthening)

Furthermore, if F is an independent set of size $\lambda_{\max}(M)$, then $\mathbb{1}_F$ lies in eigenspace of $\lambda_{\max}$

→ Yields Uniqueness

of gap between $\lambda_{\max} \neq \lambda_2 \rightarrow$ stability results

→ Cross Intersecting Families.

Theorem (Hoffman):

Let M be a $V \times V$ matrix s.t

$$\textcircled{1} M(x,y) = 0 \quad \text{if } (x,y) \notin E$$

$$\textcircled{2} M\mathbb{1} = \mathbb{1}$$

$$\text{Then, } \alpha(G) \leq \frac{-\lambda_{\min}}{1 - \lambda_{\min}} |V|$$

If F attains the bound then $\frac{\mathbb{1}_F - \frac{|E|}{|V|}\mathbb{1}}{|V|}$ lies in eigenspace of $\lambda_{\min}$

Proof: Either directly or instead by red~~u~~cing to Lovász Bound

$$\text{Let } M' = J - \frac{|V|}{1 - \lambda_{\min}} M$$

M' satisfies the cond~~i~~ of Lovász Bound.

$$J \text{ has e.val } \begin{cases} |V| & \text{w/ mult } 1 \\ 0 & \text{w/ mult } |V|-1 \end{cases}$$

$$\text{Hence, } M' \text{ has e.val } \begin{cases} \left(1 - \frac{1}{1 - \lambda_{\min}}\right) |V| & \text{w/ mult } 1 \\ \frac{-|V|}{1 - \lambda_{\min}} \lambda & \end{cases}$$

$$\text{Hence, } \lambda(M') \leq \frac{-\lambda_{\min}}{1 - \lambda_{\min}} |V|$$

Schrijver Bound

$$\theta_2(G) = \min \lambda$$

$$\begin{cases} M(x,y) \geq 1 & \text{if } (x,y) \notin E \\ M \leq \lambda I \end{cases}$$

Similar strengthening of Hoffman Bd

Back to traffic light puzzle:

$$G_n = (\mathbb{Z}_3^n, \{(x, y) \mid x_i \neq y_i, \forall i\})$$

Soln to traffic light puzzle $\rightarrow$ 3-coloring of G .

Each color class an independent set

F - one color class.

Want to apply Hoffman's bound w/
matrix $B = A/2^n$ where

$$A_n := \text{Adj}(G_n)$$

$$A_1 = \begin{pmatrix} 0 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \end{pmatrix}$$

$$A_n = A_1^{\otimes n}$$

$$A_1 \begin{pmatrix} 1 & 1 & 1 \\ 1 & \omega & \omega^2 \\ 1 & \omega^2 & \omega \end{pmatrix} = \begin{pmatrix} 2 & -1 & -1 \\ 2 & -\omega & -\omega^2 \\ 2 & -\omega^2 & -\omega \end{pmatrix} \quad \text{e.val. } \{2, -1, -1\}$$

$$\text{Eigenvalues of } A_n = 2^n, -2^{n-1}, 2^{n-2}, -2^{n-3}, \dots$$

$$\lambda_{\min}(A) = -2^{n-1}$$

—

B satisfies Hoffman cond_{ns}

$$B\mathbb{1} = \mathbb{1}$$

$$\lambda_{\min}(B) = -1/2$$

$$\therefore \text{Hence, } \Theta_H(G) = \frac{-1/2}{1 - (-1/2)} \cdot 3^n = \frac{3^n}{3} = 3^{n-1}$$

All color classes contain at most 3^{n-1} points

Each color class is exactly 3^{n-1} points

So by Hoffman Bound

$\frac{1}{F} - \frac{1}{3} \mathbb{1}$ lies in the eigenspace of $\lambda_{\min}(A)$ (ie, -2^{n-1})

Eigenspace of $-2^{n-1} = \text{span}\{\varphi_i: \mathbb{Z}_3 \rightarrow \mathbb{R} \mid \sum_{x_i} \varphi_i(x_i) = 0\}$

$$\frac{1}{F}(x_1 \dots x_n) = \frac{1}{3} + \sum_{i=1}^n \varphi_i(x_i) =: g(x_1 \dots x_n)$$

Suppose φ_1, φ_2 - non-constant

$$\left. \begin{array}{l} \varphi_1(a) < \varphi_1(b) \\ \varphi_2(c) < \varphi_2(d) \end{array} \right\} \Rightarrow \begin{array}{l} \text{Then } g(a, c, x_3 \dots x_n) \\ < g(b, c, x_3 \dots x_n) \\ < g(b, d, x_3 \dots x_n) \end{array}$$

Impossible since g attains
at most 2 possible
values.

Hence, at most one of φ_i is non-constant

$$\frac{1}{F} = \frac{1}{3} + \varphi_i(x_i) = \psi(x_i)$$

$$\text{So } F = \{x \mid x_i = j\}$$

Since all 3 color classes are disjoint, the corresponding i s are same.

Hence, traffic light depends only on one switch.

Applications of Hoffman Bound in Extremal Combinatorics

A subset $F \subseteq S_n$ is intersecting if $\forall \pi_1, \pi_2 \in F$
 $\exists i, \pi_1(i) = \pi_2(i)$

Upper Bound $F = \{ \pi \mid \pi(i) = i \}$
 $(n-1)!$

Lower Bound: Any F contains at most 1 of

$$\left. \begin{array}{ccccc} 1 & 2 & 3 & 4 & 5 \\ 2 & 3 & 4 & 5 & 1 \\ 3 & 4 & 5 & 1 & 2 \\ 4 & 5 & 1 & 2 & 3 \\ 5 & 1 & 2 & 3 & 4 \end{array} \right\} \text{ Hence } |F| \leq \frac{n!}{n} = (n-1)!$$

Thm [Ellis-Friedgut-Pipe]

For all $t \geq$ large enough n (depending on t)
the maximal size of a t -intersecting family
of permutations is $(n-t)!$

Moreover, this is achieved only on " t -double-coset."

$\therefore$ Let $q = p^m, k$

Consider $\binom{[n]}{k}_q = \{ \text{subspaces of } \mathbb{F}_q^n \text{ of dim } k \}$

A family of subspaces is t -intersecting if
intersection of any 2 subspaces has $\dim \geq t$.

Theorem [Frankl-Wilson]

$$\text{max-size} = \max \left\{ \binom{n-t}{k-t}_q, \binom{2k-t}{k}_q \right\}$$

Domain: all perfect matching in K_{2n}
(or almost p.m. in K_{2n+1})

Thm: [Lindzey 2019] Same as EFP.

Next lecture:

Domain: all subgraphs of K_n

Δ -intersecting family of graphs is
a collection of graphs, the intersection
of any two of which contains a ΔK_1 .