

Today

- Polar Codes II

CSS.318.1

Coding Theory

Lecture 22 (2022-11-16)

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Recap Polar Codes from last lecture

① Linear Compression Scheme for $(\text{Ber}(\rho))^n$

$$H: \mathcal{E}_0, \mathcal{B}^n \rightarrow \mathcal{E}_0, \mathcal{B}^m \quad m \leq (H(\rho) + \epsilon)n$$

Decompressor : D s.t

$$\Pr_{\substack{Z \sim (\text{Ber}(\rho))^n}} [D(HZ) = Z] \leq \epsilon.$$

n, R, T of D - $\text{poly}(\frac{1}{\epsilon})$

② P- invertible matrix ; $W = PZ$

Defn $|S_\epsilon| = \underbrace{\{c \in [n] / H(W_i / W_{\neq i}) \geq \epsilon\}}_{\text{unpredictable bits}}$

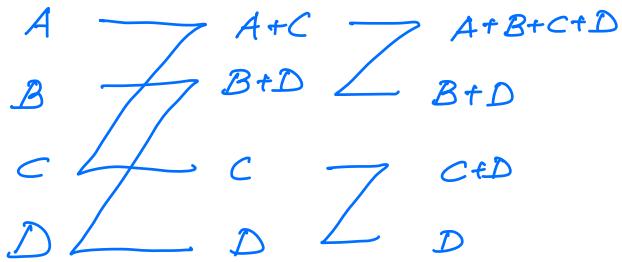
(ϵ, τ) - Polarizing matrix if $|S_\epsilon| \leq (H(\rho) + \epsilon)n$

③ Arikant's construction:

$$P_2: \mathcal{E}_0, \mathcal{B}^2 \rightarrow \mathcal{E}_0, \mathcal{B}^2$$



$$P_{2n}(U, V) = (P_n(U+V), P_n(V))$$



Defn: (ε, τ) -polarization of

$$\Pr_{i \in [n]} [H(W_i | W_{\leq i}) \in (\tau, 1-\tau)] \leq \varepsilon.$$

Lemma: (ε, τ) -polarization \Rightarrow Linear compression ($\tau \leq 1/2$)

scheme w/ $m = 151 = (H(\theta) + \varepsilon + 2\tau)n$

• Decompressor D s.t.

$$\Pr [D(HZ) + Z] \leq \tau \cdot n.$$

Theorem: For any matrix A satisfies $(\frac{1}{n^\alpha}, \frac{1}{n^\beta})$ -polarization

Consequence of Lemma + Thm: we obtain $n = \text{poly}(\frac{1}{\varepsilon})$.

$$\varepsilon = \frac{1}{n^\alpha} \quad n = \left(\frac{1}{\varepsilon}\right)^{1/\alpha} = \text{poly}\left(\frac{1}{\varepsilon}\right).$$

\rightarrow Decompressor } Both can be done

\rightarrow Find S } in $\text{poly}\left(\frac{1}{\varepsilon}\right)$ -time
(but not in today's lecture).

Proof of Lemma:

$$H = \{c \in [n] \mid H(w_c / w_{ci}) \geq 1 - \varepsilon\}$$

$$L = \{c \in [n] \mid H(w_c / w_{ci}) \leq \varepsilon\}.$$

$$M = \{c \in [n] \mid H(w_c / w_{ci}) \in (\varepsilon, 1 - \varepsilon)\}$$

$$(c, \varepsilon)\text{-polarization} \Rightarrow |M| \leq \varepsilon n$$

$$\textcircled{1} \quad |H| + |L| + |M| = n$$

$$\begin{aligned} \textcircled{2} \quad |H|(1 - \varepsilon) &\leq H(\rho)n \Rightarrow |H| \leq \frac{H(\rho)n}{1 - \varepsilon} \\ &\leq H(\rho)(1 + 2\varepsilon)n \\ &\leq (H(\rho) + 2\varepsilon)n. \end{aligned}$$

$$S \triangleq H \cup M$$

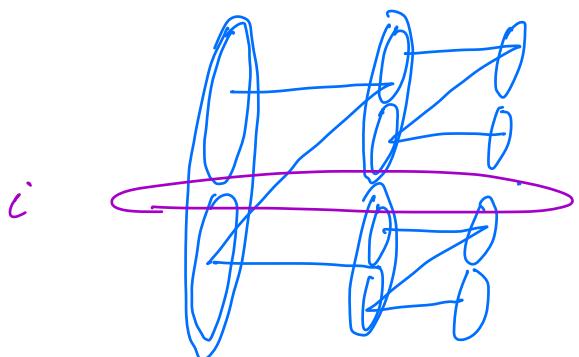
$$|S| \leq (H(\rho) + 2\varepsilon + \varepsilon)n$$

Compressor

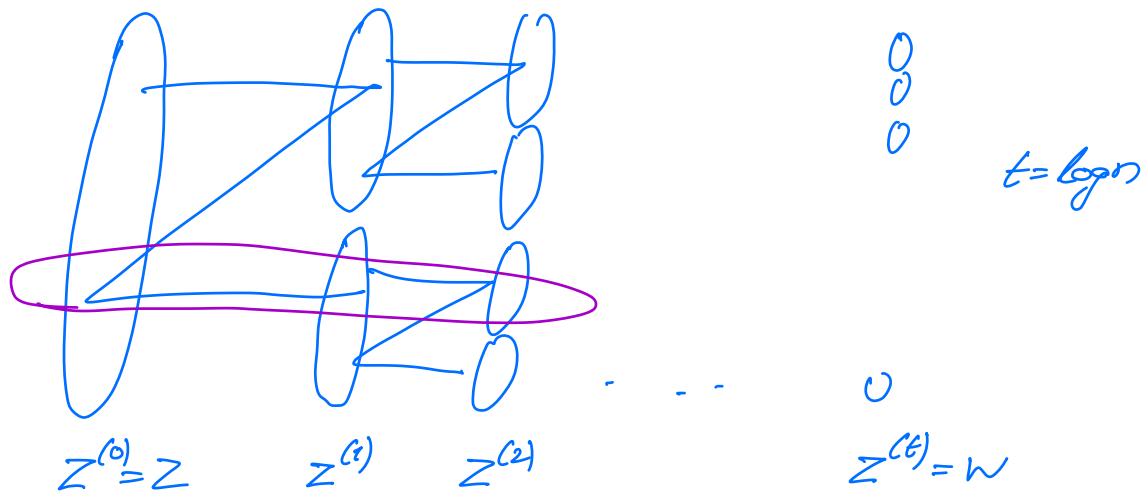
$$Z \mapsto (PZ)|_S$$

↗

Theorem: ^{KG, Fox} An $k \times k$ matrix satisfies $(\frac{1}{n^\alpha}, \frac{1}{n^\beta})$ -polarization



$$\begin{aligned} &H(w_i / w_{ji}) \\ &\text{hard to prove this.} \\ &\Pr_{c \in [n]} \left[H(w_i / w_{ji}) \in \left(\frac{1}{n^{1-\alpha}}, \frac{1}{n^{\alpha}} \right) \right] \\ &\leq \frac{1}{n^{0.001}} \end{aligned}$$



Pick $i \in [n]$, $X_i \triangleq H(Z_i^{(0)} | Z_{\leq i}^{(0)})$

X_0, X_1, \dots, X_t - Amari's martingale.

Def: X_0, X_1, \dots, X_t - martingale in M

If $t_j \in \{1, \dots, t\}$, $\forall a_0, \dots, a_{j-1} \in M$

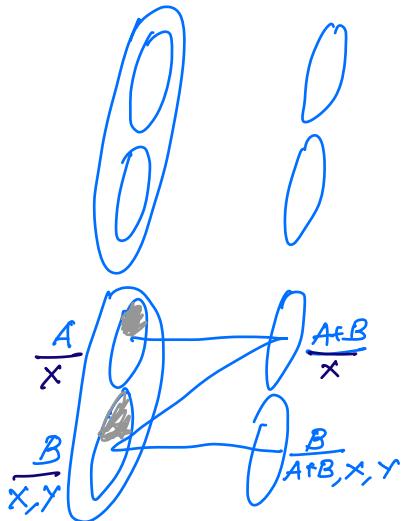
$$E[X_j | X_0 = a_0, X_1 = a_1, \dots, X_{j-1} = a_{j-1}] = a_j$$

In our setting

$$X_0 = H(p)$$

$$\begin{array}{ccc} A & \xrightarrow{\hspace{2cm}} & A+B \\ & \searrow & \swarrow \\ B & & B \end{array}$$

$$\begin{aligned} H(A, B) &= H(A+B, B) \\ H(A) + H(B|A) &= H(A+B) + H(B|A+B) \\ H(A) + H(B) &= H(A+B) + H(B|A+B) \end{aligned}$$



Initial sum

$$H(A/x) + H(B/x, y) \\ = H(A/x) + H(B/y)$$

Ending sum

$$H(A+B/x) + H(B/A+B, x, y)$$

$$H(A, B/x, y) = H(A+B, B/x, y)$$

$$\text{RHS: } H(A+B/x, y) + H(B/A+B, x, y)$$

$$\begin{aligned} \text{LHS} &= H(A/x, y) + H(B/A, x, y) \\ &= H(A/x) + H(B/y) \end{aligned}$$

Want to prove:

Strong Polarization ($t = \log n$)

$$\Pr[X_t \in (2^{-100t}, 1 - 2^{-100t})] \leq \frac{1}{2^{0.001t}}$$

Examples of martingales in [0, 1]

$$\textcircled{1} \quad X_0 = \frac{1}{2} : \quad X_t = \begin{cases} X_{t-1} + \frac{1}{2} e_t & \text{a/p } \frac{1}{2} \\ X_{t-1} - \frac{1}{2} e_t & \text{a/f } \frac{1}{2} \end{cases}$$

$$\textcircled{2} \quad X_0 = p \quad X_{t+1} = \begin{cases} X_t + \frac{1}{2} \min\{X_t, 1-X_t\} & \text{a/p } \frac{1}{2} \\ X_t - \frac{1}{2} \min\{X_t, 1-X_t\} & \text{a/f } \frac{1}{2}. \end{cases}$$

$$\textcircled{3} \quad X_0 = \alpha \quad X_{t+1} = \begin{cases} X_t^2 & \text{if } P \\ 2X_t - X_t^2 & \text{if } \bar{P}. \end{cases}$$

Local Polarization: X_0, X_1, \dots X_t -martingale satisfies local polarization.

(i) Variance in the Middle:

$$\forall c, \exists \varepsilon, \forall t$$

$$\text{Var}(X_t \mid X_{t-1} \in (c, 1-c)) \geq \sigma^2$$

(ii) Suction of the ends

$$\forall c > 1, \exists \varepsilon, \forall t$$

$$0\text{-end. : } P_t \left[X_t < \frac{X_{t-1}}{c} \mid X_{t-1} < \varepsilon \right] \geq \frac{1}{2}.$$

$$1\text{-end. : } P_t \left[1-X_t < \frac{1-X_{t-1}}{c} \mid 1-X_{t-1} < \varepsilon \right] \geq \frac{1}{2}.$$

Thm: Azéma-Karie martingale satisfies local polarization

(calculus involving entropy
will skip \approx Taylor approx)

Thm: Local Polarization \Rightarrow Strong Polarization.

(Recall Strong Polarization)

$$P_t \left[X_t \in \left(2^{-100t}, 1-2^{-100t} \right) \right] \leq \frac{1}{2^{0.001t}}$$

Proof sketch of local \Rightarrow strong.

$$\Phi_t = \min\{\sqrt{x_t}, \sqrt{1-x_t}\}$$

Claim: $\exists \beta < 1$, st $E[\Phi_t | x_{t-1}] \leq \beta \cdot \Phi_{t-1}$

(variance in the middle

section of odds).

skip proof.

$$E[\Phi_t] \leq \beta^t$$

$$P_n[\bar{\Phi}_t \geq \beta^{t/2}] \leq \beta^{t/2} \quad (\text{by Markov})$$

$$\text{i.e. } P_n[\min(x_t, 1-x_t) \geq \beta^t] \leq \beta^{t/2}.$$

$$P_n[x_t \in (\beta^t, 1-\beta^t)] \leq \beta^{t/2}$$

$$(2^{-0.0001t}, 1-2^{-0.0001t}) \leq 2^{-0.0005t}$$

not good enough *good*

Semi-strong polarization.

Two phases

(1) $t/2$ phases

obtain semi-strong
polarization

$$\frac{1}{2} \underbrace{\frac{3}{2}}_{\sim} 2^{-0.001t}$$

$$\frac{1}{2} \underbrace{\frac{3}{2}}_{\sim} 0$$

(2) Next ℓ_2 phases:

with probability e^{ℓ_2}/c , the marchmark never crosses ε .

- Get in this case at shrinks by factor c at least ℓ the time ε doubles at most ℓ the time.

□